

Key to Cryptography

I finally developed what I think is a fairly elegant solution to this Kattis Problem (it took me longer than it should, I think!)

Here's the code I wrote to solve this problem: Key to Cryptography.

```

1 #include <iostream>
2 #include <string>
3 using namespace std;
4
5 int main() {
6     string cypherText, secretWord;
7     cin >> cypherText >> secretWord;
8     string plainText = cypherText;
9     int cypherLength = cypherText.size(), secretLength = secretWord.size();
10    int minLength = cypherLength < secretLength ? cypherLength : secretLength;
11    int prt1{}, prt2{};
12    for (int i = 0; i < minLength; i++) {
13        plainText[prt1++] = (char)((((cypherText[i]-secretWord[i]+26) % 26)+65));
14    }
15
16    for (int i = secretLength; i < cypherLength; i++) {
17        plainText[prt1++] = (char)((((cypherText[i]-plainText[prt2++]+'A') % 26)+65));
18    }
19
20    cout << plainText;
21 }

```

Answer the following questions (all essay questions) about this code in the quiz posted on Canvas.

1. Explain why, on line 8, it makes sense to have `plainText` initialized with the value of `cypherText`.
2. Here we'll try to get at what's going on on line 10.
 - (a) What is the value of `minLength` if `cypherText` = "ULIXYDI" and `secretWord` = "CHIEF"?
 - (b) What is the value of `minLength` if `cypherText` = "ULIXY" and `secretWord` = "SEATTLE"?
 - (c) What's the purpose of `minLength` in the for-loop that follows in lines 12-14?
3. Suppose `cypherText` = "ULIXYDI" and `secretWord` = "CHIEF". Complete the table below to see how the first part of `plainText` is computed in the for-loop of lines 12-14.

threading for-loop				
prt1	i	cypherText[i]	secretWord[i]	(cypherText[i] - secretWord[i] + 26) % 26
0	0	U	C	18
⋮	⋮	⋮	⋮	⋮

4. Now complete the table below to see how the second part of `plainText` is computed in the for-loop of lines 16-18.

threading for-loop					
prt1	prt2	i	cypherText[i]	plainText[i]	(cypherText[i]- plainText[prt2++]+ 'A') % 26
	0				